EECS 483: Compiler Construction Lecture 10: **Dynamic Typing**

February 17 Winter Semester 2025







Announcements

- Assignment 3 released

 - Start early!

- extern functions, lambda lifting, SysVAMD64 calling convention

State of the Snake Language

- Adder: Straightline Code (arithmetic circuits) Boa: local control flow (finite automata) Cobra: procedures, extern (pushdown automata)
- Remaining limitations:
- 1. Only data are ints (booleans are really just special ints)
- 2. Only ways to use memory are local variables and the call stack





State of the Snake Language

Adder: Straightline Code (arithmetic circuits) Boa: local control flow (finite automata) Cobra: procedures, extern (pushdown automata)

Snake v4: Diamondback

- 1. Add new datatypes, use dynamic typing to distinguish them at runtime
- 2. Include heap-allocated variable-sized arrays, allowing for unrestricted memory usage Computational power: Turing complete





Booleans in Boa/Cobra

In Boa/Cobra, booleans and integers weren't truly distinct datatypes.

- All integers could be used in logical operations
- All booleans could be used in arithmetic operations

Booleans in Boa/Cobra

- -1 && 3
- true + 5
- 7 >= false

Let's change the language semantics so these are errors instead.

The following are all valid programs with well-defined semantics in Boa/Cobra

Booleans in Boa/Cobra

Can we implement operations **isInt** and **isBool** that distinguish between integers and booleans?

- isInt(true) == false
- isInt(1) == true

No, true and 1 have the exact same representation at runtime

Static vs Dynamic Typing

How would we implement a language where integers and booleans were considered disjoint?

- 1. Static Typing (C/C++, Java, Rust, OCaml) Identify the runtime types of all variables in the program Reject type-based misuse of values in the frontend of the compiler.
- 2. Dynamic Typing (JavaScript, Python, Ruby, Scheme) Use **type tags** to identify the type of data at runtime Reject type-based misuse of values at runtime, right before the operation is performed

Example 1:

true + 5

Static typing: compile time error: true used where integer expected

- Dynamic typing: runtime error: addition operation expects inputs to be integers

Example 2:

def main(x): x + 5

Static typing: need to declare a type for x, in this case int

Dynamic typing: succeed at runtime if x is an int, otherwise fail

Example 2:

def main(a): def complex_function(): ... in let x = if complex_function(): 1 else: true x + 5

Dynamic typing: succeed at runtime if **complex_function** returns true, otherwise fail

- Static typing: reject this program, even if **complex_function** always returns true

Static Typing

the runtime representation of our compiled values

Dynamic Typing

types are possible

- Easier on the compiler: if type information is reliable, we can use that to inform
- Easier on the programmer? Types document the code, aid in tooling, design

Easier on the programmer? Complex patterns that are difficult to assign static

Poll: Is static typing or dynamic typing better?

My opinion:

implementing well.

In Assignment 4, we'll implement dynamic typing

In Assignment 5, perform optimizations to reduce the runtime overhead of dynamic typing

later in the course.

- I prefer static typing, but both are popular enough to be worth studying and
- Revisit syntactic aspects of static typing and the relation with static analysis

Semantics of Dynamic Typing

Live code interpreter

Semantics of Dynamic Typing

- when we should error and how to implement isInt, isBool.
- need to specify
 - what the appropriate error messages are
 - evaluation order between expressions executing and type tags
 - true + (let $_{-} = print(3)$ in 3)
 - does this print 3 before it errors?

 A Snake value is not just an int anymore. It is either an int or a boolean, and we need to be able to tell the difference at runtime in order to determine

• Many operations can now produce runtime errors if type tags are incorrect,

Representing Dynamically Typed Values In Adder/Boa/Cobra, all runtime values were integers. In Diamondback, a runtime value must have both a type tag and a value that matches the type tag

How should we represent tags and values in our compiled program?

Representing Dynamically Typed Values Approach 1: Values as 8 bytes, Tags as extra data

A snake value is 9 bytes byte to keep our values byte-aligned

Upside: Faithful representation of our Rust interpreter pervasively

- the first byte is a tag: 0x00 for integer, 0x01 for boolean. Use a full
- the remaining 64 bits are the underlying integer, bool or pointer
- Downside: 1 byte memory overhead for all values plus padding, calling convention and architecture are 8-byte oriented, tedious to implement

Representing Dynamically Typed Values Approach 2: Values as pointers

A snake value is a 64-bit pointer to an object on the heap pointer is always 64 bits.

A value stored in this way is called **boxed**.

- value stored on the heap can then be whatever size we want, the
- store a tag and value on the heap similarly to previous approach.

Representing Dynamically Typed Values Approach 2: Values as pointers

A snake value is a 64-bit pointer to an object on the heap pointer is always 64 bits.

Downside: memory overhead. Accessing the tag requires a non-local memory access, performing an arithmetic operation multiple Approach taken in Python

- value stored on the heap can then be whatever size we want, the
- store a tag and value on the heap similarly to previous approach.
- Upside: uniform implementation, 64-bit values can be compiled as before

Representing Dynamically Typed Values Approach 3: compromise

A snake value is a 64-bit value. Use the least significant bits of the value as a **tag**. Represent simple data like integers, booleans within the 64-bits Represent large datatypes like arrays, closures, structs as pointers to the heap

Upside: use stack allocation more often Downside: can't fit 64 bits and a tag... Roughly the approach used in high-performance Javascript engines (v8) as well as some garbage-collected typed languages (OCaml)

Representing Dynamically Typed Values To implement our compiler, we need to specify 1. How each of our Snake values are represented at runtime 2. How to implement the primitive operations on these representations

Integers Implement a snake integer as a 63-bit signed integer followed by a 0 bit to indicate that the value is an integer

Number	Representation	
1	0b00000000_000000000_00000010	
6	0b00000000_000000000_00001100	
-1	0b11111111_1111111_1111110	

I.e., represent a 63-bit integer **n** as the 64-bit integer 2 * n

Booleans

The least significant bit must be 1 to distinguish from integers

datatypes

Number	Represent
true	0b000000
false	0b000000

2^62 - 2 bit patterns are therefore "junk" in this format

- Use least significant bits 0b01 to distinuish from integers and other
- Use the remaining 62 bits to encode true and false as before as 1 and 0

tation

00_0000....0000_00000101

00_0000....0000_0000001

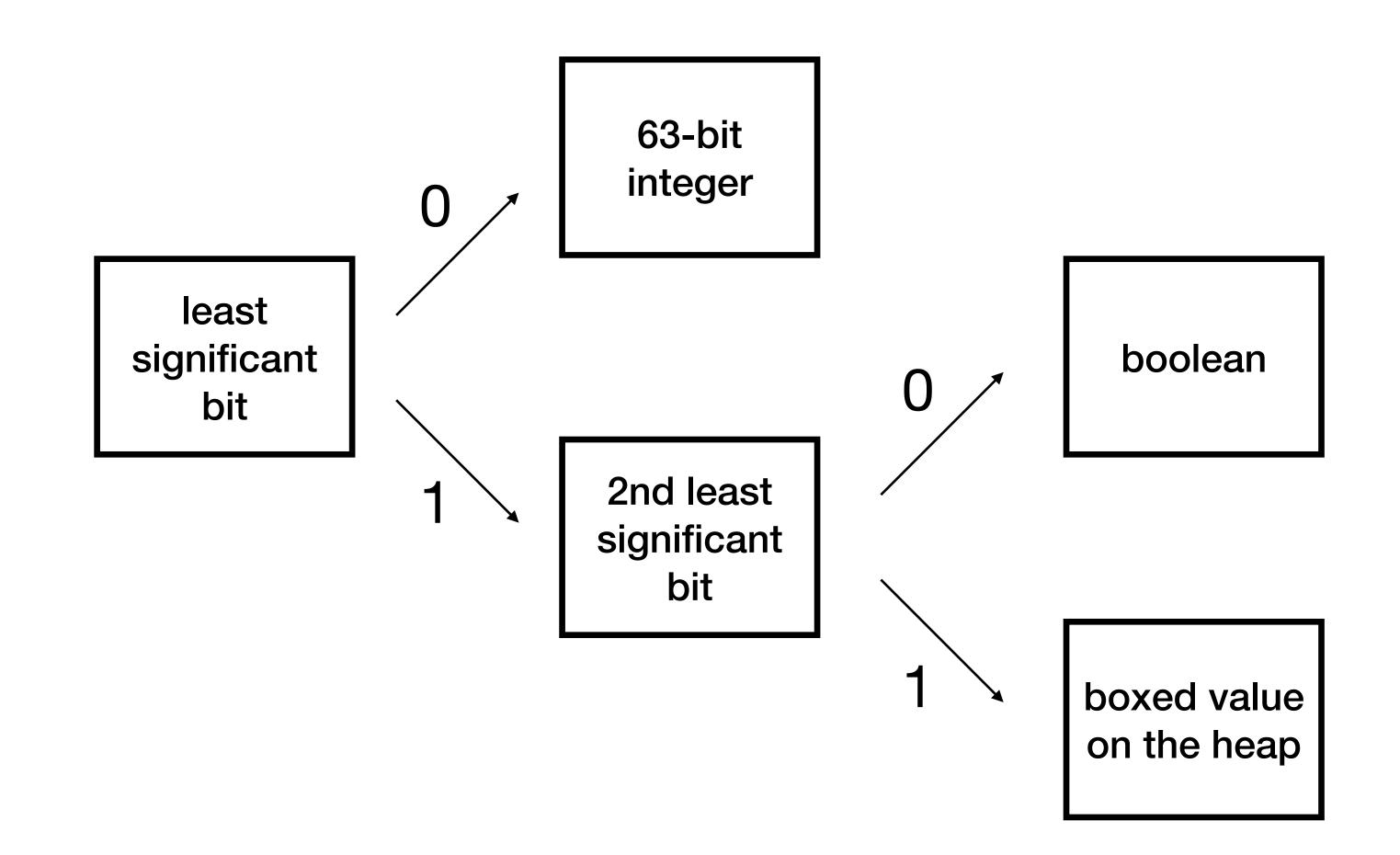


Boxed Data

The least significant bit must be 1 to distinguish from integers Use least significant bits 0b11 to distinguish from booleans.

Use remaining 62-bits to encode a pointer to the data on the heap Why is this ok? Discuss more thoroughly on Wednesday

Representing Dynamically Typed Values



Implementing Dynamically Typed Operations

- 1. Arithmetic operations (add, sub, mul)
- 2. Inequality operations (<=, <, >=, >)
- 3. Equality
- 4. Logical operations (&&, ||, !)

As well as supporting our new operations is int and is Bool

We need to revisit our implementation of all primitives in assembly code to see how they should work with our new datatype representations.



Implementing Dynamically Typed Operations In dynamic typing, implementing a primitive operation has two parts: 1. How to check that the inputs have the correct type tag 2. How to actually perform the operation on the encoded data



Implementing Dynamically Typed Operations Live code

